

AN ELEMENTARY APPROACH TO THE STACK SIZE OF REGULARLY DISTRIBUTED BINARY TREES

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For the sake of brevity we assume that the reader has a certain knowledge of [2]. Let T be a *binary tree* with n leaves. Evaluating T in postorder it is assumed that in one unit of time a node is stored in the stack or is removed from the top of the stack. Consider the number of nodes stored in the stack after t units of time. Let $R_d(n, t)$ denote the d -th moment with respect to the origin of this statistic. In [1] R. Kemp was able to produce an *exact formula* for $R_1(n, t)$ by use of 2 combinatorial identities. These identities are generalized and more easily proved in [3]. As stated in [2], a similar approach would give *exact formulas for* $R_d(n, t)$, d odd, *but not for even* d . For that purpose R. Kemp used a complex variable approach to give *asymptotic equivalents* for the numbers $R_d(n, t)$, assuming that $n, t \rightarrow \infty$ and $t \rightarrow 2\rho n$, $0 < \rho < 1$, ρ a constant. He obtains

$$(1) \quad R_d(n, t) = \pi^{-1/2} 2^{d+1} \Gamma\left(\frac{d+3}{2}\right) \cdot (n\rho(1-\rho))^{d/2} + O(n^{(d-1)/2}).$$

Here we show that it is possible to give exact formulas for $R_d(n, t)$ for all d by *elementary methods*. For instance, we give a formula for $R_2(n, t)$ from which an exact formula for the variance can be obtained by

$$\sigma^2(n, t) = R_2(n, t) - R_1^2(n, t)$$

and Kemp's formula for $R_1(n, t)$ (see [1]).

$$R_1(n, 2t) = \frac{2t(n-t)(2n-1)}{(n-1)n} \cdot \varphi(n, t)$$

$$(2) \quad R_1(n, 2t+1) = \frac{(2t+1)(2n-2t-1)-n}{n-1} \cdot \frac{2n-1}{2n-2t-1} \cdot \frac{n-t}{n} \cdot \varphi(n, t)$$

with

$$\varphi(n, t) = \binom{2t}{t} \binom{2n-2t}{n-t} \binom{2n}{n}^{-1}.$$

It is known [2] that

$$(3) \quad R_d(n, 2t+s) = \frac{2(2n-1)}{(2t+s)(2n-2t-s)} \binom{2n}{n}^{-1} \cdot \sum_{k \geq 0} (2k+s)^{d+2} \binom{2t+s}{t-k} \binom{2n-2t-s}{n-t-s-k}.$$

Let

$$(4) \quad f(d, s; m, n) := \sum_{k \geq 1} (2k+s)^d \binom{2m+s}{m-k} \binom{2n+s}{n-k},$$

$d, s, m, n \in N_0$. We propose to show how a closed formula for $f(d, s; m, n)$ can be obtained which is obviously equivalent to the same problem for $R_d(n, t)$. The method is essentially included in [3].

Theorem 1. The following recursion holds for the numbers $f(d, s; m, n)$:

$$(5) \quad f(d+2, s; m, n) = (2m+s)^2 f(d, s; m, n) - 4(2m+s)_2 f(d, s; m-1, n).$$

Here, $(x)_k$ denotes the *falling factorials*.

Proof. Since

$$(6) \quad \binom{2m-2+s}{m-1-k} = \frac{(m-k)(m+k+s)}{(2m+s)_2} \binom{2m+s}{m-k}$$

and

$$(7) \quad 4(m-k)(m+k+s) = (2m+s)^2 - (2k+s)^2,$$

a rearrangement of (6) and summation over $k \geq 1$ gives (5). \square

So if we have formulas for $d = 0$ and 1 , we have solved our problem. $f(1, s; m, n)$ is known [3]:

$$(8) \quad f(1, s; m, n) = \binom{2m+s}{m} \binom{2n+s}{n} \frac{mn}{m+n+s}.$$

Theorem 2.

$$(9) \quad f(0, s; m, n) = \frac{1}{2} \left[\binom{2m+2n+2s}{m+n+s} - \sum_{0 \leq k \leq s} \binom{2m+s}{m+k} \binom{2n+s}{n+k} \right].$$

Proof. Let without loss of generality $m \leq n$.

$$\begin{aligned} \xi &= \sum_{0 \leq k \leq 2m+s} \binom{2m+s}{k} \binom{2n+s}{m+n+s-k} + \binom{2m+s}{m} \binom{2n+s}{n} \\ &\quad - \sum_{m \leq k \leq 2m+s} \binom{2m+s}{k} \binom{2n+s}{n-m+k} \\ &= \binom{2m+2n+2s}{m+n+s} + \binom{2m+s}{m} \binom{2n+s}{n} \\ &\quad - \xi - \sum_{1 \leq k \leq s} \binom{2m+s}{m+k} \binom{2n+s}{n+k}, \end{aligned}$$

which gives (9), since

$$\begin{aligned} \xi &= \sum_{0 \leq k \leq m} \binom{2m+s}{m-k} \binom{2n+s}{n-k} \\ &= f(0, s; m, n) + \binom{2m+s}{m} \binom{2n+s}{n}. \quad \square \end{aligned}$$

Obviously, (9) is only useful for small values of s . However, in practice we require just $s = 0$ and $s = 1$. By doing some computations using Theorem 1 we obtain

Corollary 3.

$$(10) \quad f(2, 0; m, n) = \binom{2m+2n}{m+n} \frac{2mn}{2m+2n-1},$$

$$(11) \quad f(4, 0; m, n) = \binom{2m+2n}{m+n} \frac{8mn(3mn-m-n)}{(2m+2n-1)(2m+2n-3)}.$$

$$\begin{aligned} &f(2, 1; m, n) \\ &= \frac{1}{2} \binom{2m+2n+2}{m+n+1} \frac{2m+2n+1+4mn}{2m+2n+1} \\ (12) \quad &\quad - \binom{2m+1}{m} \binom{2n+1}{n} \\ &= f(0, 1; m, n) + \frac{2mn}{2m+2n+1} \binom{2m+2n+2}{m+n+1}, \end{aligned}$$

$$\begin{aligned}
& f(4, 1; m, n) \\
&= \frac{1}{2} \binom{2m+2n+2}{m+n+1} \frac{(2m+1)(2n+1)(12mn+2m+2n-1)}{(2m+2n+1)(2m+2n-1)} \\
(13) \quad & - \binom{2m+1}{m} \binom{2n+1}{n} \\
&= f(2, 1; m, n) \\
&+ \binom{2m+2n+2}{m+n+1} \frac{6mn(2m+1)(2n+1)}{(2m+2n+1)(2m+2n-1)}. \quad \square
\end{aligned}$$

Since

$$\begin{aligned}
& R_d(n, 2t+s) \\
&= \frac{2(2n-1)}{(2t+s)(2n-2t-s)} \binom{2n}{n}^{-1} \\
(14) \quad & \cdot \left[s^{d+2} \binom{2t+s}{t} \binom{2n-2t-s}{n-t} \right. \\
& \left. + f(d+2, s; t, n-t-s) \right],
\end{aligned}$$

an obvious computation gives

Theorem 4.

$$\begin{aligned}
& R_2(n, 2t) = \frac{12t(n-t) - 4n}{2n-3}, \\
(15) \quad & R_2(n, 2t+1) = \frac{12t(n-t-1) + 2n-3}{2n-3}. \quad \square
\end{aligned}$$

It is trivial to obtain an asymptotic formula for $R_2(n, t)$, $n \rightarrow \infty$, $t \rightarrow \rho n$.

It is worthwhile to discuss the limitations of our elementary approach: A general asymptotic formula like (1) *cannot* be obtained, although the recursion (5) seems to be promising. But unfortunately the two terms of the right-hand side of (5) are of equal rate of growth.

REFERENCES

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